

Rank of Sparse Bernoulli Matrices

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Singularity of discrete random matrices

Conjecture 1: Let B be the $n \times n$ matrix random sign matrix.

$$\mathbb{P}(s_n(B) = 0) = \left(\frac{1}{2} + o(1)\right)^n.$$

For $n \times n$ matrix A , $s_1(A) \geq s_2(A) \geq \dots \geq s_n(A) \geq 0$ denote its singular values in non-increasing order.

$$s_k(A) = \max_{k\text{-dim } H} \min_{x \in H \cap S^{n-1}} \|Ax\|.$$

What it means? Roughly 2 columns or 2 rows are identical up to the sign difference.

Singularity of discrete random matrices

1967	Komlós:	$o_n(1)$
1995	Kahm, Komlós, and Szemerédi	$\exp(-cn)$
2006, 2007	Tao and Vu	$\left(\frac{3}{4} + o(1)\right)^n$
2010	Bourgain, Vu, and Wood	$\left(\frac{1}{\sqrt{2}} + o(1)\right)^n$
2018	Tikhomirov	$\left(\frac{1}{2} + o(1)\right)^n$

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Conjecture 2: Let A be the $n \times n$ random matrix with i.i.d Bernoulli(p) entries, where $p = p_n \in (0, \frac{1}{2})$. Then,
 $\mathbb{P}(s_n(A) = 0) = (1 + o(1))\mathbb{P}(\text{either a row or column of } A \text{ equals } 0)$

- The case when $p < \log n / n$ is immediate. The probability that A has a 0 row of column is a high probability event.

Singularity of discrete random matrices

Theorem 3 (Litvak-Tikhomirov): Let A be the $n \times n$ random matrix with i.i.d Bernoulli(p) entries, where $p = p_n$ satisfies

$$C \frac{\log(n)}{n} \leq p \leq C^{-1}.$$

Then,

$$\begin{aligned} & \mathbb{P}(s_n(A) \leq t) \\ & \leq (1 + o_n(1)) \mathbb{P}(\text{either a row or a column of } A \text{ equals } 0) + \exp(3 \log(n)^2) t \end{aligned}$$

If $p \geq q$ for some constant $q > 0$.

$$\begin{aligned} & \mathbb{P}(s_n(A) \leq t) \\ & \leq (1 + o_n(1)) \mathbb{P}(\text{either a row or a column of } A \text{ equals } 0) + C_q n^{2.5} t \end{aligned}$$

Singularity of discrete random matrices

- Jain, Sah, and Sawhney settle the case when $c < p < \frac{1}{2}$.
- Basak and Rudelson settle the case when $\frac{\log(n)}{n} \leq p \leq (1 + o_n(1)) \frac{\log(n)}{n}$.



Rank of discrete random matrices

Conjecture 4: Let B be the $n \times n$ matrix random sign matrix. For any non-negative integer k ,

$$\mathbb{P}(s_{n-k}(B) = 0) = \left(\frac{1}{2} + o(1)\right)^{kn}.$$

Rank of discrete random matrices

Theorem (Costello, Vu)

Let A_{sym} be a n by n symmetric Bernoulli matrix for $\frac{(1+\epsilon) \log n}{2n} \leq p < c$. Then, with high probability, the corank of A_{sym} is the number of zero columns/rows. (c can be greater than $\frac{1}{2}$. Also some extensions from the A_{sym} model)

Corollary (Basak-Rudelson, Litvak-Tikhomirov, Jain-Sah-Sawhney)

For $p \in \left[\frac{\log n}{n}, \frac{(1+o(1)) \log n}{n} \right) \cup \left(\frac{C \log n}{n}, \frac{1}{2} \right)$, with high probability, the corank of A is the maximum of number of zero columns and number of zero rows.

Rank of sparse Bernoulli matrices

Theorem 5 (H.): Let A be the $n \times n$ random matrix with i.i.d Bernoulli(p) entries, where $p = p_n$ satisfies

$$1 \leq \liminf \frac{pn}{\log(n)} \leq \limsup \frac{pn}{\log(n)} < +\infty.$$

Then, for any positive integer k ,

$$\begin{aligned} & \mathbb{P}(s_{n-k+1}(A) \leq t) \\ & \leq (1 + o_n(1)) \mathbb{P}(\text{either } k \text{ rows or } k \text{ columns of } A \text{ equal } 0) + n^{2k+o(1)}t. \end{aligned}$$

- Setting $t=0$, a similar type of result for Conjecture 4, but for sparse Bernoulli matrices.
- When $k = 1$, settle Conjecture 2 (singularity of Bernoulli matrices) in the regime of p described above.

Magnitude of the event

Let $O_{RC}(k)$ be the event of A that either k rows or k columns of A are zero.

$$\mathbb{P}(O_{RC}) = O(\exp(-(r-1)k \log n)),$$

where $r = \frac{pn}{\log n} \geq 1$.

Definition A probability event O of A is called a (r, k) –high probability event if

$$\mathbb{P}(O^c) = o(\mathbb{P}(O_{RC}(k))).$$

Decomposition of \mathbb{R}^n from Litvak-Tikhomirov

For $x \in \mathbb{R}^n$, x^* be the *non-increasing rearrangement* of x .

Let $\sigma_x: [n] \mapsto [n]$ be the permutation so that $|x_{\sigma_x(i)}| \geq |x_{\sigma_x(j)}|$ whenever $i \geq j$.

$$x_i^* := |x_{\sigma_x(i)}|.$$

Decomposition of \mathbb{R}^n from Litvak-Tikhomirov

- \mathcal{V} -vectors: collection of vectors x in \mathbb{R}^n so that $i \mapsto \frac{x^*(i)}{x^*(\lambda n)}$ is not growing too fast as $i \searrow 1$.

Decomposition of \mathbb{R}^n from Litvak-Tikhomirov

- \mathcal{R} -vectors: collection of vectors x such that there exists $1 \leq n_1 \leq n_2 \leq n$ such that $\frac{\|x_{[n_1, n_2]}^*\|_\infty}{\|x_{[n_1, n_2]}^*\|_2} \leq \frac{c}{p}$.

Decomposition of \mathbb{R}^n from Litvak-Tikhomirov

- \mathcal{J} -vectors: collection of vectors x such that $x_{n_1}^* \gg x_{n_2}^*$ for $(n_1, n_2) \in S$.

Structure of the Proof in Litvak-Tikhomirov

- $\{\forall x \in \mathcal{V}, \|Ax\| \gg 0\}$ is a $(r, 1)$ - high probability event.
- $\{\forall x \in \mathcal{R}, \|Ax\| \gg 0\}$ is a $(r, 1)$ - high probability event.
- $\{\forall x \in \mathcal{T}, \|Ax\|, \|A^\top x\| \gg 0\}$ is a $(r, 1)$ - high probability event if we condition on $O_{RC}^c(1)$.
- $\mathcal{V} \cup \mathcal{R} \cup \mathcal{T} = \mathbb{R}^n$.

Structure of the Proof

With some modifications of \mathcal{V} , \mathcal{R} , and \mathcal{T} ,

- $\{\forall x \in \mathcal{V} \cap \mathbb{R}_J, \|A_{I_1, J_1} x\| \gg 0\}$ is a (r, k) - high probability event.
- $\{\forall x \in \mathcal{R}, \|Ax\| \gg 0\}$ is a (r, k) - high probability event.
- $\{\exists I_1, J_1 \subset [n]$ with $|I| = |J| = n - k + 1$ such that
 $\forall x \in \mathcal{T} \cap \mathbb{R}_J, \|A_{I_1, J_1} x\| \gg 0$, and $\forall x \in \mathcal{T} \cap \mathbb{R}_I, \|A_{I_1, J_1}^\top x\| \gg 0\}$ is a (r, k) - high probability event if we condition on $O_{RC}^c(k)$.

\mathcal{T} -vectors

Condition on typical Column support sizes

- $\mathcal{L}(t)$ be the collection of columns whose support has size at most t .
- $J_0 = \mathcal{L}(t)$ for t depends on r, k .
- $I_0 = \{i \in [n], \exists j \in \mathcal{L}(t) \text{ s.t. } a_{ij} = 1\}$ (union of supports of columns from J)

Matrix Decomposition:

When $(r - 1)k < 1 - o(1)$

When $(r - 1)k < 1 - o(1)$

When $(r - 1)k \geq 1 - o(1)$

When $(r - 1)k \geq 1 - o(1)$