

# On the $\ell_0$ Isoperimetry of Measurable Sets

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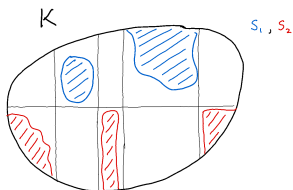
AGA Seminar

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# Axis-disjointness

## Definition (Axis-Disjointness)

Let  $S_1, S_2 \subseteq \mathbb{R}^n$ . We say  $S_1, S_2$  are axis-disjoint if there does not exist a line parallel to some coordinate axis that intersects both  $S_1$  and  $S_2$ .



## $\ell_0$ Isoperimetric inequalities

Suppose  $K \subseteq \mathbb{R}^n$  is convex and  $S_1, S_2 \subseteq K$  are axis-disjoint. and  $S_3 = K \setminus (S_1 \cup S_2)$

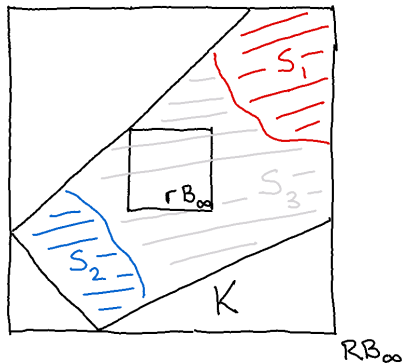
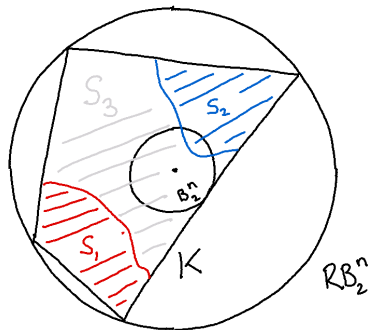
- Laddha, Vempala [2]: If  $B_2 \subseteq K \subseteq RB_2$  and  $\varepsilon > 0$  then

$$\mathbf{vol}(S_3) \geq \Omega \left( \frac{\varepsilon}{n^{3.5}R} \cdot (\min\{\mathbf{vol}(S_1), \mathbf{vol}(S_2)\} - \varepsilon \cdot \mathbf{vol}(K)) \right)$$

- Narayanam, Rajaraman, Srivastava [3]: If  $rB_\infty \subseteq K \subseteq RB_\infty$  then

$$\mathbf{vol}(S_3) \geq \Omega \left( \frac{(r/R)}{n^{3.5}} \min\{\mathbf{vol}(S_1), \mathbf{vol}(S_2)\} \right)$$

# Depiction of $\ell_0$ isoperimetry



# $\ell_0$ Isoperimetric coefficient

## Definition ( $\ell_0$ Isoperimetric coefficient)

Let  $K \subseteq \mathbb{R}^n$  be measurable. The  $\ell_0$  isoperimetric coefficient of  $K, \psi_K$ , is defined as

$$\psi_K := \inf_{\substack{S_1, S_2 \subseteq K \\ S_1, S_2 \text{ are measurable} \\ S_1, S_2 \text{ are axis disjoint}}} \frac{\text{vol}(K \setminus (S_1 \cup S_2))}{\min\{\text{vol}(S_1), \text{vol}(S_2)\}}$$

## Theorem (Laddha, Vempala [2])

Let  $C$  denote any axis-aligned cube. Then

$$\psi_C \geq \frac{\log 2}{n}.$$

Proof is via the Loomis-Whitney inequality.

# Main result: improved $\ell_0$ isoperimetry

## Theorem (F.)

Let  $\mathcal{C}$  denote any axis-aligned cube. Then there exists absolute constants  $c_1, c_2 > 0$  such that

$$\frac{c_1}{\sqrt{n}} \leq \psi_{\mathcal{C}} \leq \frac{c_2}{\sqrt{n}}.$$

## Corollary (F. Improved $\ell_0$ isoperimetric inequalities)

- If  $B_2 \subseteq K \subseteq RB_2$  and  $\varepsilon > 0$  then

$$\mathbf{vol}(S_3) \geq \Omega \left( \frac{\varepsilon}{n^3 R} \cdot (\min\{\mathbf{vol}(S_1), \mathbf{vol}(S_2)\}) - \varepsilon \cdot \mathbf{vol}(K) \right)$$

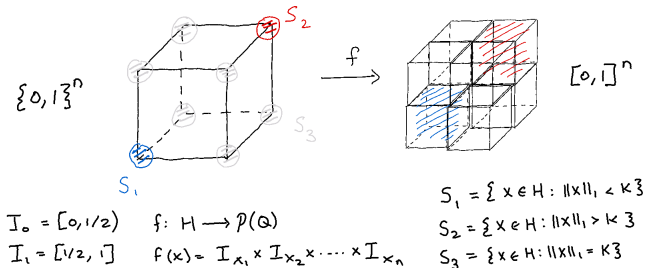
- If  $rB_\infty \subseteq K \subseteq RB_\infty$  then

$$\mathbf{vol}(S_3) \geq \Omega \left( \frac{(r/R)}{n^3} \min\{\mathbf{vol}(S_1), \mathbf{vol}(S_2)\} \right)$$

# Upper bound: cube case

## Lemma

Let  $K \subseteq \mathbb{R}^n$  be an axis-aligned cube. Then  $\psi_K = O(n^{-1/2})$

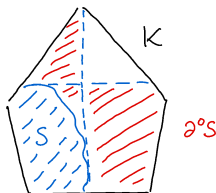


observation:  $f(S_1), f(S_2)$  are axis-disjoint!

# $\ell_0$ boundary

## Definition ( $\ell_0$ boundary)

Let  $K \subseteq \mathbb{R}^n$  be measurable and  $S \subseteq K$  be measurable. The  $\ell_0$  boundary of  $S$ ,  $\partial^0 S$  is defined as the set of all points in  $K \setminus S$  that are not axis-disjoint from  $S$ .



## Remark

$$\psi_K = \inf_{\substack{S_1, S_2 \subseteq K \\ S_1, S_2 \text{ are measurable} \\ S_1, S_2 \text{ are axis disjoint}}} \frac{\text{vol}(\partial^0 S_1)}{\min\{\text{vol}(S_1), \text{vol}(S_2)\}}.$$



# Structured subsets of $[0, 1]^n$

## Definition

Let  $A \subseteq [0, 1]^n$ .

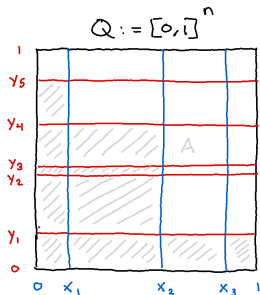
- 1 We say that  $A$  is anchored if for every  $x \in A$  the rectangle  $[0, x_1] \times \cdots \times [0, x_n]$  is contained in  $A$ .
- 2 We say  $A$  has a grid structure if there exists a grid such that  $A$  can be written as a union of grid blocks.

## Definition

Given  $r \in \{0, 1, \dots, n\}$  and  $p \in [0, 1]$  we define the  $p$ -weighted hamming ball of radius  $r$ ,  $H(p, r)$ , as

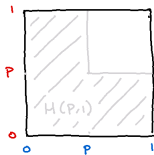
$$H(p, r) := \{x \in [0, 1]^n : \#\{i : x_i \leq p\} \geq n - r\}.$$

# Harper's Theorem for structured subsets [1]



$\mathcal{A} := \{A \subseteq Q \text{ such that}$   
 ①  $A$  has a grid structure  
 ②  $A$  is anchored

$$H(p, r) := \{x \in Q : \#\{i : 0 \leq x_i \leq p\} \geq n - r\}$$



Thm: Harper

Fix  $0 < \lambda < 1$ . Then  $\min_{\substack{A \in \mathcal{A} \\ \text{vol}(A) = \lambda}} \text{vol}(\partial^* A)$  is achieved by  $H(p, r)$   
 for some  $0 < p < 1$ ,  $0 < r < n$

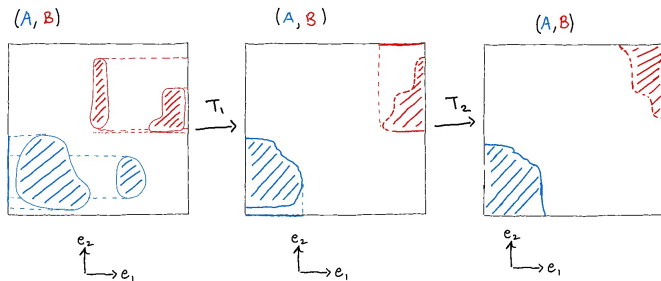
Goal: Reduce to Harper

# Continuous compression via shaking

$$(A, B) \xrightarrow{T_i} (\hat{A}, \hat{B})$$

$$\hat{A} := \bigcup_{x \in \text{Proj}_{e_i^\top} A} \{x\} \times [0, \text{vol}_1(\{y \in A \mid \text{Proj}_{e_i^\top} y = x\})]$$

$$\hat{B} := \bigcup_{x \in \text{Proj}_{e_i^\top} A} \{x\} \times (1 - \text{vol}_1(\{y \in B \mid \text{Proj}_{e_i^\top} y = x\}), 1)$$



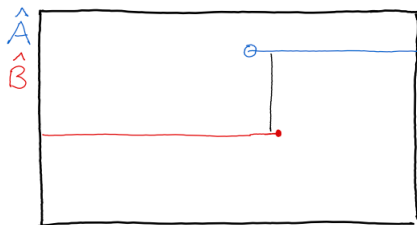
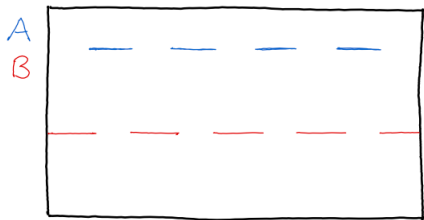
# Compression guarantees

## Lemma

If  $A, B \subseteq [0, 1]^n$  then  $(\hat{A}, \hat{B})$  satisfies the following:

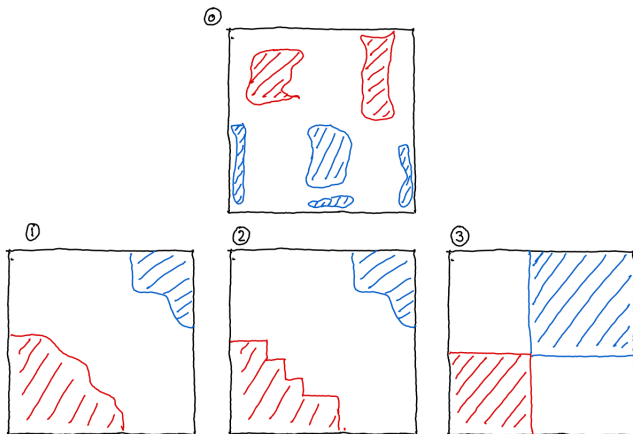
- 1  $\mathbf{vol}(A) = \mathbf{vol}(\hat{A}), \mathbf{vol}(B) = \mathbf{vol}(\hat{B})$ .
- 2 If  $A, B$  are axis-disjoint then  $\hat{A}, \hat{B}$  are axis-disjoint.
- 3 The boundary of  $\hat{A}$  has measure 0.

Illustration of (2)



# General proof strategy

Overview: 1) Compress 2) Under approximate 3) Harper



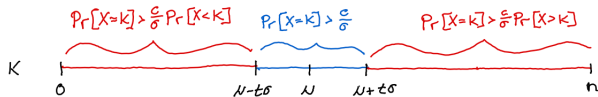
# Binomial inequality

Observation: The volume of  $B(p, r)$  is equal to the probability that a binomial random variable,  $\text{Bin}(n, 1 - p)$ , is at most  $r$ .

## Lemma

There exists a universal constant  $c > 0$  such that the following is true: Let  $X \sim \text{Bin}(n, 1 - p)$  be a binomial random variable. Let  $k \in \{0, 1, \dots, n\}$ . Then the following inequality holds:

$$\mathbb{P}[X = k] \geq \frac{c}{\sqrt{np(1-p)}} \min \{ \mathbb{P}[X < k], \mathbb{P}[X > k] \}$$



$$\begin{aligned} X &\sim \text{Binom}(n, 1-p), \quad p \leq 1/2 \\ \mu &= n(1-p) \\ \sigma &= \sqrt{np(1-p)} \end{aligned}$$

## Theorem (F. General upper bound)

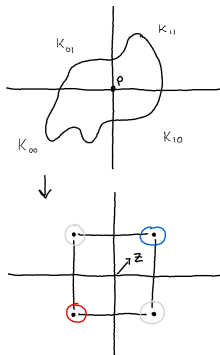
*There exists a universal constant  $c > 0$  such that the following holds: Let  $K \subseteq \mathbb{R}^n$  be a measurable set. Then*

$$\psi_K \leq cn^{-1/2}$$

- To prove the upper bound we develop a general framework for constructing axis-disjoint subsets  $S_1, S_2 \subseteq K$  for which  $\frac{\text{vol}(K \setminus (S_1 \cup S_2))}{\min\{\text{vol}(S_1), \text{vol}(S_2)\}}$  is small.
- Our proof is an application of the Probabilistic Method.

# Random splitting plane

Construction



1) Pick  $x \in \mathbb{R}^n$  such that

$$\text{vol}(K \cap \{x_i \leq p_i\}) = \text{vol}(K \cap \{x_i > p_i\})$$

2) Decompose  $K$  into orthants induced by  $P$

3) Pick  $z$  uniformly from  $\{0, 1\}^n$

4) Define

$$S_1 := \bigcup_{\|s-z\|_1 < L/2} K_s \quad S_2 := \bigcup_{\|s-z\|_1 > \Gamma/2} K_s$$

$$S_3 := \bigcup_{\|s-z\|_1 \in \{L/2, \Gamma/2\}} K_s$$

Goal: Show that with non-zero probability

$$\text{vol}(S_3) = O(n^{-1/2}), \text{vol}(S_1), \text{vol}(S_2) = \Omega(1).$$



# Symmetric case

Let  $K \subseteq \mathbb{R}^n$  be symmetric:

- First moment argument: With any constant probability  $\mathbf{vol}(S_3) = O(n^{-1/2})$
- symmetric property:  $\mathbf{vol}(S_1) = \mathbf{vol}(S_2)$

## Lemma

Let  $K \subseteq \mathbb{R}^n$  be symmetric. Then for every  $z \in \{0, 1\}^n$  we have

$$\sum_{\|s-z\|_1 < \lfloor n/2 \rfloor} \mathbf{vol}(K_s) = \sum_{\|s-z\|_1 > \lceil n/2 \rceil} \mathbf{vol}(K_s).$$

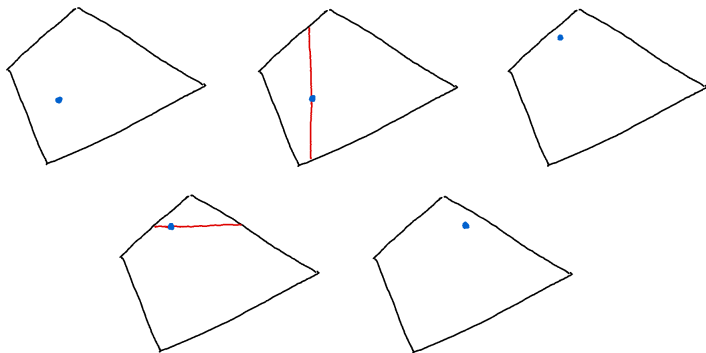
Consequence:  $\mathbf{vol}(S_1) = \mathbf{vol}(S_2) = \Omega(1)$ .

# Non-symmetric case

- Problem: If  $K$  is non-symmetric then  $\mathbf{vol}(S_1), \mathbf{vol}(S_2)$  need not be equal.
- Workaround: second moment method
- Need to show:  $\mathbb{E}[\mathbf{vol}(S_1)^2] < 1/2 - c$  for some absolute constant  $c$ .
- Idea 1: If  $s_1, s_2$  are far apart then  $K_{s_1}, K_{s_2}$  are more likely to be on opposite sides.
- Idea 2: The volume of  $K$  is not too concentrated amongst a cluster of orthants.

# Random Walk: Coordinate Hit-and-Run

- Introduced by Turchin (1971) [4]



- Simple to implement, good in practice.
- No strong theoretical guarantees until recently.

# Problem: lowerbounding conductance

Theorem (Lovasz, Simonivits)

$$1) \chi^2(Q, Q_t) \leq \chi^2(Q_0, Q) (1 - \phi^2/2)^t$$

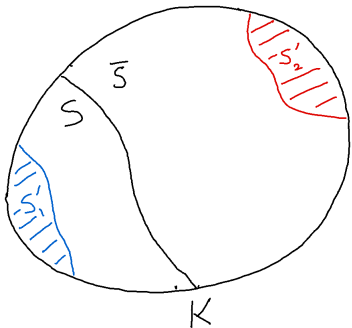
$$2) d_{TV}(Q, Q_t) \leq sM + M(1 - Q_s/2)^t$$

$$\phi(s) := \frac{\int_{x \in S} P_x(\bar{S}) dQ}{\min(Q(S), Q(\bar{S}))}$$

$$\phi := \inf_{S \subseteq K} \phi(s)$$

$$\phi_s := \inf_{\substack{S \subseteq K \\ \text{vol}(S) \leq 1/2}} \phi(s)$$

$$M = \sup_{S \subseteq K} \frac{Q(S)}{Q(S_0)}$$



$$S'_1 = \left\{ x \in S : P_x(\bar{S}) < \frac{1}{2n} \right\}$$

$$S'_2 = \left\{ x \in \bar{S} : P_x(S) < \frac{1}{2n} \right\}$$

Lemma (Laddha, Vempala)

$S'_1, S'_2$  are axis disjoint

# Open questions

- What is the best lowerbound on the isoperimetric coefficient for convex bodies  $K$ ? Is it also  $\Omega(n^{-1/2})$ ?
- Can the upper bound on  $\psi_K$  be improved by depending on the regularity of  $K$ ?

Thank you!



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